

# A leakage-resilient MAC

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## Side-channel attacks

- Side-channel attacks

- Leakage-resilient cryptography
- Leakage-resilient authentication
- Stateful MACs
- A leakage-resilient stream cipher [Pie09]
- Tree-based leakage-resilient stream cipher
- Questions?
- References

Side-channel attacks (SCAs) are attacks that exploit (physical) properties of the *implementation*, e.g. power use. Even the best smart cards and similar devices are vulnerable to SCAs.

- Practitioners have tried to solve this
  - Ad-hoc
  - Only partially succesful
- Leakage-resilient cryptography is the theoretical approach
  - Inspired by provable security
  - Requires a good model of what SCAs can do

# Leakage-resilient cryptography

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We use the model of continuous leakage proposed by Dziembowski and Pietrzak [DP08, Pie09].

- Mostly equal to the standard model...
- ...but the adversary can supply a *leakage function* with each input, and receives the output of this function with the output.
  - this function must produce  $\lambda$  bits of output, where  $\lambda$  depends on the (quality of) implementation
  - only computation leaks information
  - this function must be efficient

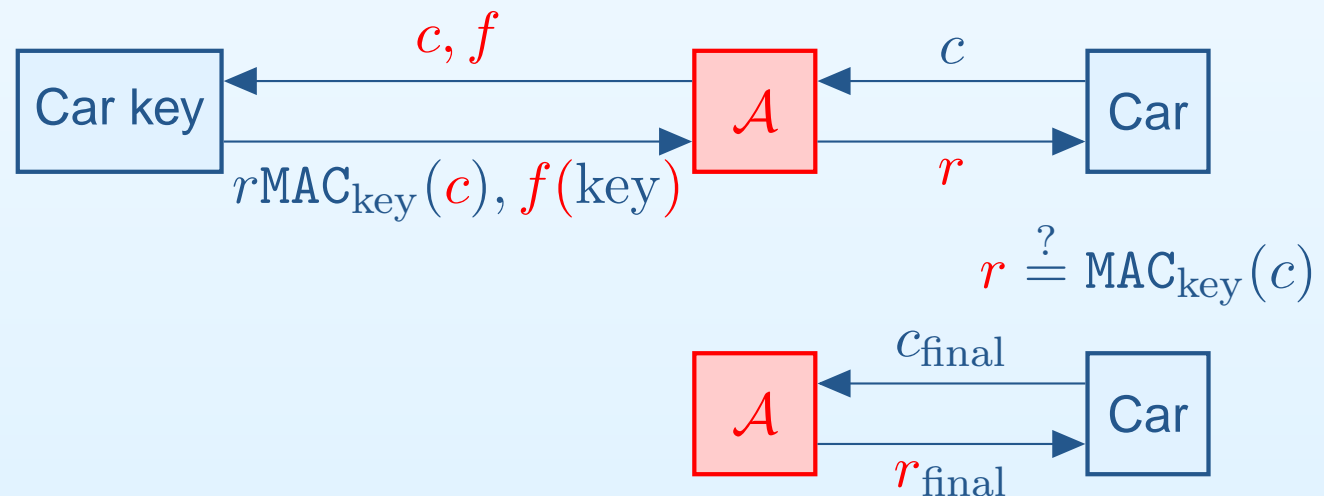
# Leakage-resilient authentication

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Inspired by the power analysis of KeeLoq [EKM<sup>+</sup>08].



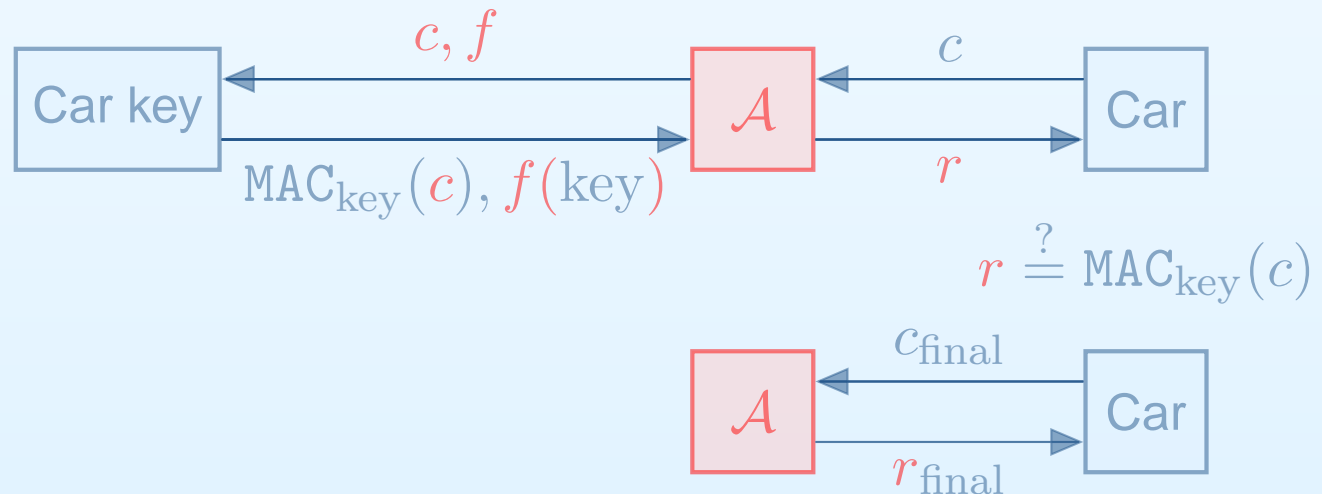
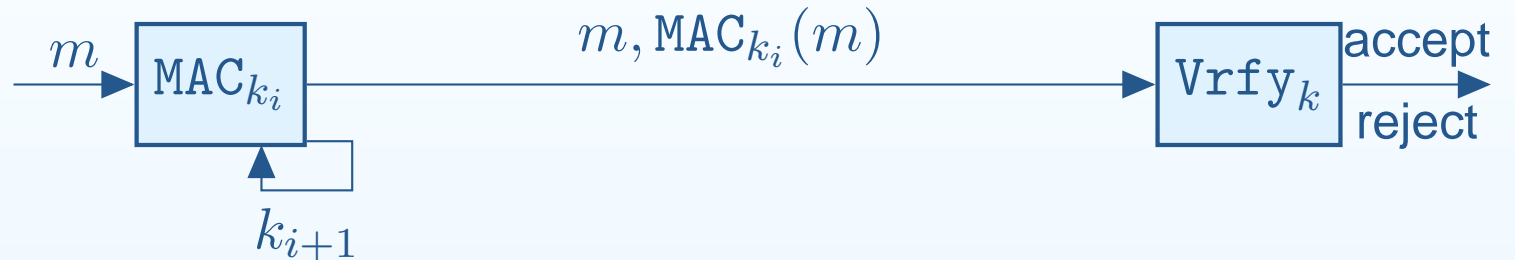
The adversary has temporary access and wins if  $r_{\text{final}}$  is valid. The classical solution is a MAC. Of course, we need leakage-resilience.



# Stateful MACs

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Problem: if we use the same key each time, it will eventually completely leak. So we need a *stateful* MAC.

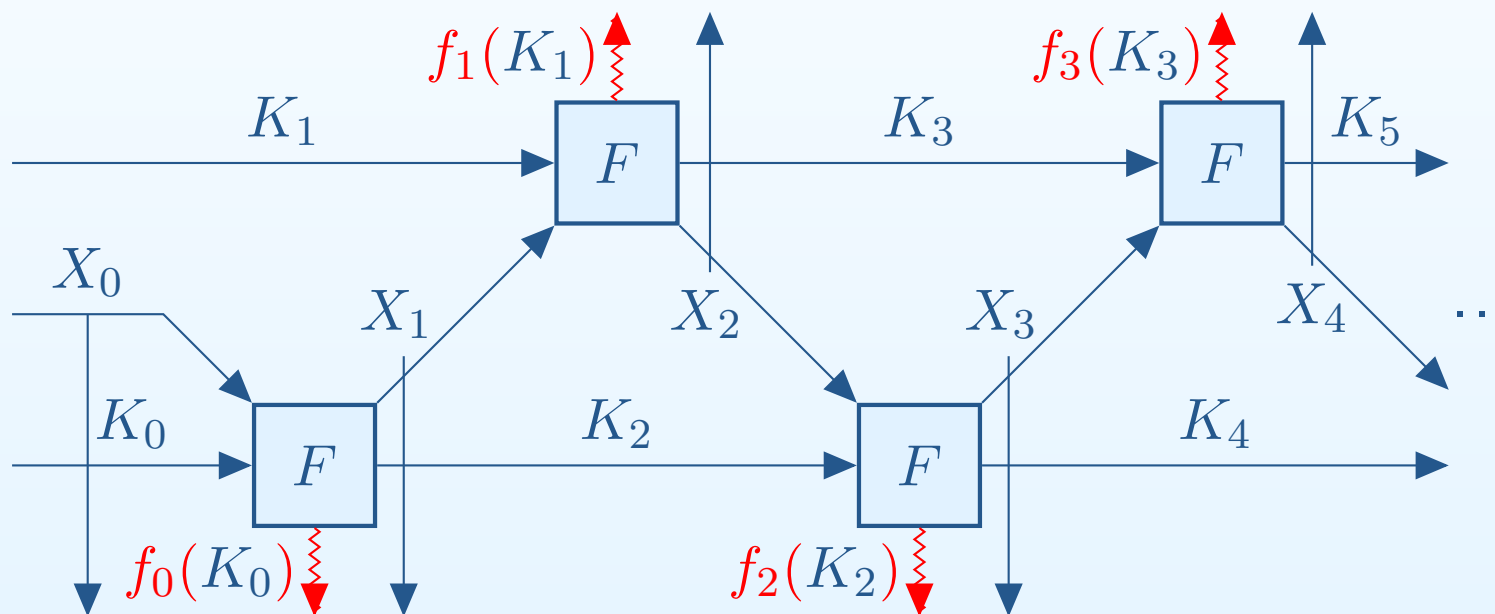


## A leakage-resilient stream cipher [Pie09]

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Stream cipher:  $X_i$  is pseudorandom given  $X_0, X_1, \dots, X_{i-1}$ .

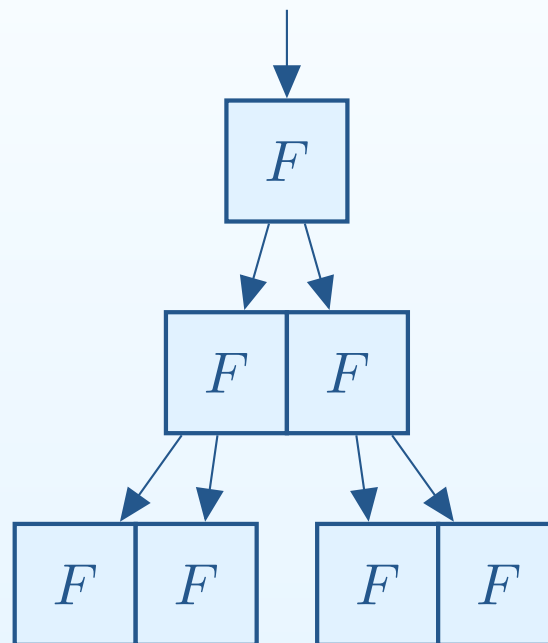
Leakage-resilient stream cipher:  $X_i$  is pseudorandom given  $X_0, X_1, X_2, \dots, X_{i-1}$  and the leakage  $f_0(K_0^+), f_1(K_1^+), \dots, f_{i-1}(K_{i-1}^+)$  from these rounds.



# Tree-based leakage-resilient stream cipher

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We are working on this.



The authenticating side performs a depth-first search on the tree. The verifier only needs to perform  $O(\log(\#\text{queries}))$  calculations to calculate any output.

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- **References**

## References

- [DP08] Stefan Dziembowski and Krzysztof Pietrzak. Leakage-resilient cryptography. In *FOCS*, pages 293–302. IEEE Computer Society, 2008.
- [EKM<sup>+</sup>08] Thomas Eisenbarth, Timo Kasper, Amir Moradi, Christof Paar, Mahmoud Salmasizadeh, and Mohammad T. Manzuri Shalmani. On the power of power analysis in the real world: A complete break of the KeeLoq code hopping scheme. In David Wagner, editor, *CRYPTO*, volume 5157 of *Lecture Notes in Computer Science*, pages 203–220. Springer, 2008.
- [Pie09] Krzysztof Pietrzak. A leakage-resilient mode of operation. In Antoine Joux, editor, *EUROCRYPT*, volume 5479 of *Lecture Notes in Computer Science*, pages 462–482. Springer, 2009.